

**1. Introduction.** This program translates input from `/dev/kbd` to Chinese runes. The idea is based on `ktrans` but this is a rewrite to have much cleaner code and the efficiency for handling a much bigger dictionary.

The usage is the almost the same as `ktrans`: just invoke `cim` before `rio`.

`^T` gives you English mode, `^N` enters Cangjie mode. `^L` rotates selection.

I know you might want Pinyin but Cangjie is a little bit easier to implement and way easier to use when there is no prompt.

**2. Setup keyboard I/O.** See `kbdfs(8)`. Reading `/dev/kbd` gives keycodes and the `c` messages are what we want to process. Besides handling the event produced by the operating system, we need also to ensure other programs that read `/dev/kbd` can acquire the translated result by faking a new keyboard device to shadow the original one.

```

⟨ Setup keyboard 2 ⟩ ≡
    if ((kbd = open("/dev/kbd", OREAD)) < 0)
        sysfatal("open_kbd:_%r");
    ⟨ Shadowing /dev/kbd 4 ⟩

```

This code is used in section 30.

**3.** ⟨ Variables 3 ⟩ ≡  
**int** *kbd*, *newkbd*;

See also section 6.

This code is used in section 30.

**4.** The call `bind("#|", dir, MREPL)` would create a pair of pipes named `data` and `data1` in `dir`.

```

⟨ Shadowing /dev/kbd 4 ⟩ ≡
    if (bind("#|", "/n/temp", MREPL) < 0)
        sysfatal("bind_/n/temp:_%r");
    if ((newkbd = open("/n/temp/data1", OWRITE)) < 0)
        sysfatal("open_kbd_pipe:_%r");
    if (bind("/n/temp/data", "/dev/kbd", MREPL) < 0)
        sysfatal("bind_kbd_pipe:_%r");
    unmount(nil, "/n/temp");

```

This code is used in section 2.

**5.** This loop would be used to intercept input and output translated result. When the start of the message is `'c'` the following string contains one character. The function `nextstate()` would output runes to `newkbd` when possible.

```

⟨ Key translation loop 5 ⟩ ≡
    while ((n = read(kbd, buf, sizeof (buf))) > 0) {
        buf[n - 1] = 0;
        if (n < 2 ∨ buf[0] ≠ 'c')
            write(newkbd, buf, n);
        else nextstate(buf, n, newkbd);
    }

```

This code is used in section 30.

**6.** ⟨ Variables 3 ⟩ +≡  
**int** *n*;  
**char** *buf*[128];

**7. Dictionary.** The dictionary is defined as several C arrays in a separated file named `dict.c`. It has to be generated by a external program and that detail would not be discussed here. The requirement for dictionary is that all entries need to be sorted with lexicographical order. An `int` is used as a key to encode the string of at most three alphabets.

```
format Rune char /* Rune is the type for Unicode characters */
⟨Declarations 7⟩ ≡
typedef struct Map Map;
struct Map {
    int key;
    Rune *val;
};
extern const Map dict[];
```

See also section 8.

This code is used in section 30.

**8.** There are two precomputed indices to speed up lookup.

```
⟨Declarations 7⟩ +≡
extern const int wl1[];
extern const int wl2[];
```

**9.** `wl1` contains indices from "a" to "z".

```
⟨Lookup in wl1 9⟩ ≡
return &dict[wl1[key[0] - 'a']];
```

This code is used in section 11.

**10.** `wl2` contains indices from "aa", "ab" to "zz". Since it is likely the indices are sparse, a negative value indicates not found.

```
⟨Lookup in wl2 10⟩ ≡
{
    int index = wl2[(key[0] - 'a') * 26 + (key[1] - 'a')];
    return (index < 0) ? nil : &dict[index];
}
```

This code is used in section 11.

**11.** Find in the dictionary. Since the indices are very sparse after level 2, only linear search would be needed. The optional argument `table` can be used to resume a search.

```
⟨Subroutines 11⟩ ≡
const Map *match(const char *key, int len, const Map *table)
{
    switch (len) {
    case 1: ⟨Lookup in wl1 9⟩
    case 2: ⟨Lookup in wl2 10⟩
    default:
        if (table ≠ nil) ⟨Linear search 13⟩
        else ⟨Speedup with wl2 12⟩
    }
}
```

See also section 16.

This code is used in section 30.

**12.**  $\langle$ Speedup with *lvl2* 12 $\rangle \equiv$   
**return** *match*(*key*, *len*, *match*(*key*, 2, *nil*));

This code is used in section 11.

**13.** Since the keys are ordered, if the compare result is negative, return *nil*. Only codes after the first two characters are stored in the key. If the key's length is less than 3 it is only stored as 0.

$\langle$ Linear search 13 $\rangle \equiv$   
 {  
   **int** *cmp*;  
   **int** *keycode*;  
    $\langle$ Compute keycode 14 $\rangle$   
   **for**  $\langle$ First none 0 key until next 0 15 $\rangle$  {  
     *cmp* = (*table*[*i*].*key*  $\gg$  (5 \* (3 - (*len* - 2)))) - *keycode*;  
     **if** (*cmp*  $\equiv$  0) **return** &*table*[*i*];  
     **else if** (*cmp* > 0) **return** *nil*;  
   }  
   **return** *nil*;  
 }

This code is used in section 11.

**14.** Each alphabet occupies 5 bits. There are 3 alphabets in the keycode.

$\langle$ Compute keycode 14 $\rangle \equiv$   
   *keycode* = 0;  
   **for** (**int** *i* = 2; *i* < *len*; *i*++) {  
     *keycode*  $\ll$ = 5;  
     *keycode* += *key*[*i*] - 'a' + 1;  
   }

This code is used in section 13.

**15.** When a two character key doesn't exist, the index for next key with same prefix is stored.

$\langle$ First none 0 key until next 0 15 $\rangle \equiv$   
   (**int** *i* = (*table*[0].*key*) ? 0 : 1; *table*[*i*].*key*; *i*++)

This code is used in section 13.

**16. State machine.** Right now this is a bit of messy.

```

⟨Subroutines 11⟩ +≡
void nextstate(char *oldbuf, int n, int outfd)
{
    if (len > 5) {
        len = 0;
        laststate = nil;
    }
    if (¬natural ∧ islower(oldbuf[1])) {
        ⟨Copy input character to input buffer 26⟩
        ⟨Search and output 19⟩
        return;
    }
    else ⟨Handle mode change and selection 17⟩
    ⟨Directly write to output 28⟩
    return;
}

```

**17.** ⟨Handle mode change and selection 17⟩ ≡

```

if (oldbuf[1] ≡ 20) { /* ^T */
    natural = 1;
    ⟨Reset state 25⟩
    return;
}
else if (oldbuf[1] ≡ 14) { /* ^N */
    natural = 0;
    ⟨Reset state 25⟩
    return;
}
else if (¬natural ∧ oldbuf[1] ≡ '␣') {
    ⟨Reset state 25⟩
    return;
}
else if (oldbuf[1] ≡ 12) { /* ^L */
    ⟨Rotate selection 20⟩
    return;
}
else if (oldbuf[1] ≡ '\b') {
    ⟨Reset state 25⟩
}
}

```

This code is used in section 16.

**18.** This indicates the English/Chinese state.

```

⟨Global states 18⟩ ≡
int natural = 0;

```

See also sections 21 and 24.

This code is used in section 30.

19. The only **goto** statement in this program. It enables auto start of next input if the last character is already complete.

```

⟨Search and output 19⟩ ≡
    const Map *result;
Search: result = match(input, len, laststate);
    if (result ≠ nil) {
        if (laststate ≠ nil) ⟨Backspace 23⟩
            candidate = result-val;
        ⟨Write current candidate 22⟩
        laststate = result;
    }
    else if (laststate ≠ nil) {
        ⟨Flush input buffer until the but last character 27⟩
        laststate = nil;
        goto Search;
    }

```

This code is used in section 16.

```

20. ⟨Rotate selection 20⟩ ≡
    if (laststate ≠ nil) {
        ⟨Backspace 23⟩
        candidate++;
        if (¬*candidate) candidate = laststate-val;
        ⟨Write current candidate 22⟩
    }

```

This code is used in section 17.

```

21. ⟨Global states 18⟩ +≡
    Rune *candidate;

```

```

22. ⟨Write current candidate 22⟩ ≡
    char buf[128];
    int n = snprintf(buf, sizeof (buf), "%C", *candidate) + 1;
    write(outfd, buf, n);

```

This code is used in sections 19 and 20.

```

23. ⟨Backspace 23⟩ ≡
    write(outfd, "c\b\0", 3);

```

This code is used in sections 19 and 20.

24. These states keeps a record of current input.

```

⟨Global states 18⟩ +≡
    char input[20];
    int len = 0;
    const Map *laststate = nil;

```

```

25. ⟨Reset state 25⟩ ≡
    len = 0;
    laststate = nil;

```

This code is used in section 17.

**26.** ⟨Copy input character to input buffer 26⟩ ≡  
*input*[*len*] = *oldbuf*[1];  
*len* += 1;  
*input*[*len*] = '\0';

This code is used in section 16.

**27.** ⟨Flush input buffer until the but last character 27⟩ ≡  
*input*[0] = *input*[*len* - 1];  
*input*[1] = '\0';  
*len* = 1;

This code is used in section 19.

**28.** ⟨Directly write to output 28⟩ ≡  
*write*(*outfd*, *oldbuf*, *n*);

This code is used in section 16.

**29. Program entry.** A standard Plan9 C program would use these headers.

```
#include <u.h>
#include <libc.h>
#include <ctype.h>
```

**30.** The `main()` function starts the daemon process and call `exits(nil)`.

```
<Declarations 7>
<Global states 18>
<Subroutines 11>
void main(int argc, char **argv)
{
  <Variables 3>;
  if (argc > 1) {
    <Test dictionary 31>
    exits(nil);
  }
  <Setup keyboard 2>
  if (fork()) exits(nil);
  <Key translation loop 5>
}
```

**31.** <Test dictionary 31> ≡

```
const Map *result = match(argv[1], strlen(argv[1]), nil);
if (result == nil) print("not found!\n");
else print("%S\n", result->val);
```

This code is used in section 30.



**32. Bugs.** The translation could fail if typing is too fast.

*argc*: [30](#).  
*argv*: [30](#), [31](#).  
*bind*: [4](#).  
*buf*: [5](#), [6](#), [22](#).  
*candidate*: [19](#), [20](#), [21](#), [22](#).  
*cim*: [1](#).  
*cmp*: [13](#).  
*dict*: [7](#), [9](#), [10](#).  
*dir*: [4](#).  
*exits*: [30](#).  
*fork*: [30](#).  
*i*: [14](#), [15](#).  
*index*: [10](#).  
*input*: [19](#), [24](#), [26](#), [27](#).  
*islower*: [16](#).  
*kbd*: [2](#), [3](#), [5](#).  
*kbdfs* 8: [2](#).  
*key*: [7](#), [9](#), [10](#), [11](#), [12](#), [13](#), [14](#), [15](#).  
*keycode*: [13](#), [14](#).  
*ktrans*: [1](#).  
*laststate*: [16](#), [19](#), [20](#), [24](#), [25](#).  
*len*: [11](#), [12](#), [13](#), [14](#), [16](#), [19](#), [24](#), [25](#), [26](#), [27](#).  
*lwl1*: [8](#), [9](#).  
*lwl2*: [8](#), [10](#).  
*main*: [30](#).  
**Map**: [7](#), [11](#), [19](#), [24](#), [31](#).  
*match*: [11](#), [12](#), [19](#), [31](#).  
**MREPL**: [4](#).  
*n*: [6](#), [16](#), [22](#).  
*natural*: [16](#), [17](#), [18](#).  
*newkbd*: [3](#), [4](#), [5](#).  
*nextstate*: [5](#), [16](#).  
*nil*: [4](#), [10](#), [11](#), [12](#), [13](#), [16](#), [19](#), [20](#), [24](#), [25](#), [30](#), [31](#).  
*oldbuf*: [16](#), [17](#), [26](#), [28](#).  
*open*: [2](#), [4](#).  
**OREAD**: [2](#).  
*outfd*: [16](#), [22](#), [23](#), [28](#).  
**OWRITE**: [4](#).  
*print*: [31](#).  
*read*: [5](#).  
*result*: [19](#), [31](#).  
*rio*: [1](#).  
**Rune**: [7](#), [21](#).  
*Search*: [19](#).  
*snprint*: [22](#).  
*strlen*: [31](#).  
*sysfatal*: [2](#), [4](#).  
*table*: [11](#), [13](#), [15](#).  
*unmount*: [4](#).  
*val*: [7](#), [19](#), [20](#), [31](#).  
*write*: [5](#), [22](#), [23](#), [28](#).

- ⟨ Backspace 23 ⟩ Used in sections 19 and 20.
- ⟨ Compute keycode 14 ⟩ Used in section 13.
- ⟨ Copy input character to input buffer 26 ⟩ Used in section 16.
- ⟨ Declarations 7, 8 ⟩ Used in section 30.
- ⟨ Directly write to output 28 ⟩ Used in section 16.
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- ⟨ Shadowing /dev/kbd 4 ⟩ Used in section 2.
- ⟨ Speedup with *lvl2* 12 ⟩ Used in section 11.
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- ⟨ Variables 3, 6 ⟩ Used in section 30.
- ⟨ Write current candidate 22 ⟩ Used in sections 19 and 20.

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